## Essential States in Time Petri Nets

Louchka Popova-Zeugmann

Humboldt-Universität zu Berlin Department of Computer Science

November 27, 2013

< ロ > < 同 > < 回 > < 回 >

# Outline



Time Petri Nets

- Rounding of Runs
- Essential States
- Reachable Graph









æ

ヘロン 人間 とくほど 人ほど





æ

ヘロン 人間 とくほど 人ほど





3 / 44

æ





æ





æ





3 / 44

æ





3 / 44

æ



The local time of t is reset to zero!



э

イロト イポト イヨト イヨト



The local time of *t* is reset to zero!  $\leftarrow$  static conflict

э

イロト イポト イヨト イヨト





Э

Statics:

# Petri Net (Skeleton)





э

(日)



•  $m_0 = (2, 0, 1)$ 

Э

<ロ> <同> <同> < 回> < 回>



•  $m_0 = (2, 0, 1)$  *p*-marking

Э

<ロ> <同> <同> < 同> < 同>



•  $m_0 = (2, 0, 1)$  *p*-marking •  $h_0 = (\sharp, 0, 0, 0)$  *t*-marking

<ロ> <同> <同> < 同> < 三> < 三>



•  $m_0 = (2, 0, 1)$  *p*-marking •  $h_0 = (\sharp, 0, 0, 0)$  *t*-marking

h(t) is the time shown by the clock of t since the last enabling of t

4 / 44

< D > < P > < P >

### State

The pair z = (m, h) is called a **state** in a TPN  $\mathcal{Z}$ , iff:

- m is a p-marking in  $\mathcal{Z}$ .
- *h* is a *t*-marking in  $\mathcal{Z}$ .

< ロ > < 同 > < 回 > < 回 >

Dynamics:



Let  $\mathcal{Z}$  be a TPN and let z = (m, h), z' = (m', h') be two states.  $\mathcal{Z}$  changes from state z = (m, h) into the state z' = (m', h') by:

Notation: 
$$z \xrightarrow{t} z'$$
  $z \xrightarrow{\tau} z$ 



イロト イポト イヨト イヨト



E

(ロ) (部) (目) (日)





E

ヘロト ヘ部ト ヘヨト ヘヨト



э

イロト イポト イヨト イヨト



E

ヘロン ヘロン ヘヨン ヘヨン



э

イロト イポト イヨト イヨト



(日)

#### **Definitions:**

- transition sequence:  $\sigma = t_1 \cdots t_n$
- run:  $\sigma(\tau) = \tau_0 t_1 \tau_1 \cdots \tau_{n-1} t_n \tau_n, \qquad \tau_i \in \mathbb{R}^+_0$
- feasible run:  $z_0 \xrightarrow{\tau_0} z_0^* \xrightarrow{t_1} z_1 \xrightarrow{\tau_1} z_1^* \cdots \xrightarrow{t_n} z_n \xrightarrow{\tau_n} z_n^*$
- feasible transition sequence :  $\sigma$  is feasible if there ex. a feasible run  $\sigma(\tau)$

・ロット (雪) (日) (日)

## Reachable state, Reachable marking, State space

#### **Definitions:**

- *z* is a **reachable state** in  $\mathcal{Z}$  if there ex. a feasible run  $\sigma(\tau)$ and  $z_0 \xrightarrow{\sigma(\tau)} z$
- *m* is a reachable *p*-marking in Z if there ex. a reachable state z in Z with z = (m, h)
- The set of all reachable states in Z is the state space of Z (denoted: StSp(Z)).

#### Some Problems: The State Space



The set of all reachable states is dense.

イロト イポト イヨト イヨト

# Some Further Problems: Reachability of *p*-markings



 $\mathcal{R}_{\mathcal{Z}}$  is the set if all reachable *p*-markings in Z.

 $\mathcal{R}_{\mathcal{S}(\mathcal{Z})}$  is the set of all reachable markings in the skeleton of Z ( the state space of the skeleton of Z).

・ コ ト ・ 雪 ト ・ 日 ト ・

# Some Further Problems: Reachability of *p*-markings



 $\mathcal{R}_{\mathcal{Z}}$  is the set if all reachable *p*-markings in Z.

 $\mathcal{R}_{\mathcal{S}(\mathcal{Z})}$  is the set of all reachable markings in the skeleton of Z ( the state space of the skeleton of Z).

# Some Further Problems: Reachability of *p*-markings



 $\mathcal{R}_{\mathcal{Z}}$  is the set if all reachable *p*-markings in Z.

 $\mathcal{R}_{\mathcal{S}(\mathcal{Z})}$  is the set of all reachable markings in the skeleton of Z ( the state space of the skeleton of Z).

A B A B A
A
B
A
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A

# Some Further Problems: Reachability of *p*-markings



 $\mathcal{R}_{\mathcal{Z}}$  is the set if all reachable *p*-markings in Z.

 $\mathcal{R}_{\mathcal{S}(\mathcal{Z})}$  is the set of all reachable markings in the skeleton of Z ( the state space of the skeleton of Z).

A B A B A
A
B
A
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A

# Some Further Problems: Reachability of *p*-markings



 $\mathcal{R}_{\mathcal{Z}}$  is the set if all reachable *p*-markings in Z.

 $\mathcal{R}_{\mathcal{S}(\mathcal{Z})}$  is the set of all reachable markings in the skeleton of Z ( the state space of the skeleton of Z).

A B A B A
A
B
A
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
B
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A
A

# Parametric Run, Parametric State

Let  $\mathcal{Z} = (P, T, F, V, m_0, I)$  be a TPN and  $\sigma = t_1 \cdots t_n$  be a transition sequence in  $\mathcal{Z}$ .

 $(\sigma(x), B_{\sigma})$  is a **parametric run** of  $\sigma$  and  $(z_{\sigma}, B_{\sigma})$  is a **parametric state** in  $\mathcal{Z}$  with  $z_{\sigma} = (m_{\sigma}, h_{\sigma})$ , if

- $m_0 \xrightarrow{\sigma} m_\sigma$
- $h_{\sigma}(t)$  is a sum of variables, ( $h_{\sigma}$  is a parametric *t*-marking)
- $B_{\sigma}$  is a set of conditions (a system of inequalities)

・ロト ・ 同ト ・ ヨト ・ ヨト

# Parametric Run, Parametric State

Let  $\mathcal{Z} = (P, T, F, V, m_0, I)$  be a TPN and  $\sigma = t_1 \cdots t_n$  be a transition sequence in  $\mathcal{Z}$ .

 $(\sigma(x), B_{\sigma})$  is a **parametric run** of  $\sigma$  and  $(z_{\sigma}, B_{\sigma})$  is a **parametric state** in  $\mathcal{Z}$  with  $z_{\sigma} = (m_{\sigma}, h_{\sigma})$ , if

- $m_0 \xrightarrow{\sigma} m_\sigma$
- $h_{\sigma}(t)$  is a sum of variables, ( $h_{\sigma}$  is a parametric *t*-marking)
- $B_{\sigma}$  is a set of conditions (a system of inequalities)

 $h_{\sigma}(t)$  is a **term** and  $B_{\sigma}$  is a set of **formulas** in a predicate logic (Presburger Arithmetic - decidable !)

・ロット (雪) (日) (日)
# Parametric Run, Parametric State

Let  $\mathcal{Z} = (P, T, F, V, m_0, I)$  be a TPN and  $\sigma = t_1 \cdots t_n$  be a transition sequence in  $\mathcal{Z}$ .

 $(\sigma(x), B_{\sigma})$  is a **parametric run** of  $\sigma$  and  $(z_{\sigma}, B_{\sigma})$  is a **parametric state** in  $\mathcal{Z}$  with  $z_{\sigma} = (m_{\sigma}, h_{\sigma})$ , if

- $m_0 \xrightarrow{\sigma} m_\sigma$
- $h_{\sigma}(t)$  is a sum of variables, ( $h_{\sigma}$  is a parametric *t*-marking)
- $B_{\sigma}$  is a set of conditions (a system of inequalities)

Obviously

• 
$$z_0 \xrightarrow{\sigma(x)} (z_\sigma, B_\sigma),$$

・ロト ・ 雪 ト ・ ヨ ト ・ ヨ ト

# Parametric Run, Parametric State

Let  $\mathcal{Z} = (P, T, F, V, m_0, I)$  be a TPN and  $\sigma = t_1 \cdots t_n$  be a transition sequence in  $\mathcal{Z}$ .

 $(\sigma(x), B_{\sigma})$  is a **parametric run** of  $\sigma$  and  $(z_{\sigma}, B_{\sigma})$  is a **parametric state** in  $\mathcal{Z}$  with  $z_{\sigma} = (m_{\sigma}, h_{\sigma})$ , if

- $m_0 \xrightarrow{\sigma} m_\sigma$
- $h_{\sigma}(t)$  is a sum of variables, ( $h_{\sigma}$  is a parametric *t*-marking)
- $B_{\sigma}$  is a set of conditions (a system of inequalities)

Obviously

• 
$$z_0 \xrightarrow{\sigma(x)} (z_{\sigma}, B_{\sigma}),$$
  
•  $StSp(\mathcal{Z}) = \bigcup_{\sigma, \beta} \{ z_{\sigma(\beta(x))} \mid \beta : X \to \mathbb{R}^+_0, \beta(x) \text{ satisfies } B_{\sigma} \}.$ 



ヘロト 人間 とくほ とくほ と



 $\sigma = t_4 t_3$ 





 $\sigma = t_4 t_3 \qquad : \qquad x_0$ 





 $\sigma = t_4 t_3 \qquad : \qquad x_0 t_4$ 







 $\sigma = t_4 \ t_3 \qquad : \qquad x_0 \ t_4 \ x_1$ 





 $\sigma = t_4 \ t_3 \qquad : \qquad x_0 \ t_4 \ x_1 \ t_3$ 

÷.

ヘロト ヘ団ト ヘヨト ヘヨト



 $\sigma = t_4 t_3 \qquad : \qquad x_0 t_4 x_1 t_3 x_2$ 

$$(z_{t_4t_3}, B_{t_4t_3}) = \left( \left( \begin{pmatrix} 0 \\ 1 \\ 1 \end{pmatrix}, \begin{pmatrix} x_0 + x_1 + x_2 \\ \sharp \\ x_2 \end{pmatrix} \right), \begin{cases} 2 \le x_0 \le 3, \ x_0 + x_1 \le 5, \\ 2 \le x_1 \le 4, \\ 0 \le x_2 \le 3, \\ x_0 + x_1 + x_2 \le 5 \end{cases} \right).$$



$$(\underbrace{\begin{pmatrix} 0\\1\\1 \end{pmatrix}, \begin{pmatrix} x_0 + x_1 + x_2\\ \sharp\\ x_2 \end{pmatrix}}_{z_{\sigma}}), \underbrace{\begin{pmatrix} z_{\sigma}, B_{\sigma} \end{pmatrix} = (z_{t_4 t_3}, B_{t_4 t_3}) = \\ \underbrace{\{ 2 \le x_0 \le 3, \\ \{ 2 \le x_1 \le 4, \quad x_0 + x_1 + x_2 \le 5 \} \\ 0 \le x_2 \le 3 \\ B_{\sigma} \end{bmatrix}}_{B_{\sigma}}$$

#### Bounds for the number of equalities in $B_{\sigma}$ :

Let  $\mathcal{Z} = (P, T, F, V, m_0, I)$  be a Petri net. Furthermore let  $(z_{\sigma}, B_{\sigma})$  be a parametric state and  $\sigma = t_1 \cdots t_n$  a transition sequence in  $\mathcal{Z}$ .

- The number of variables appearing in  $B_{\sigma}$  is at most n + 1.
- The number of non-redundant inequalities in  $B_{\sigma}$  is at most

$$\min\{2 \cdot (n \cdot |T| + 1), (n + 1) \cdot (\frac{n}{2} + 2)\}\$$



 $\sigma = t_1 t_3 t_4 t_2 t_3$ 

÷.



 $\sigma = t_1 t_3 t_4 t_2 t_3$ 

 $z_0 \xrightarrow{0.7}$ 

÷.



 $\sigma = t_1 t_3 t_4 t_2 t_3$ 

 $z_0 \xrightarrow{0.7}$ 

÷.



 $\sigma = t_1 t_3 t_4 t_2 t_3$ 

 $Z_0 \xrightarrow{0.7} \xrightarrow{t_1}$ 

÷.

ヘロト ヘ団ト ヘヨト ヘヨト



 $\sigma = t_1 t_3 t_4 t_2 t_3$ 

$$Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0}$$

÷.



 $\sigma = t_1 t_3 t_4 t_2 t_3$ 

$$z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0}$$

÷.



 $\sigma = t_1 t_3 t_4 t_2 t_3$ 

 $z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3}$ 

÷.

ヘロト ヘヨト ヘヨト ヘヨト



 $\sigma = t_1 t_3 t_4 t_2 t_3$ 

 $Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{0.4}$ 

÷.

ヘロト ヘヨト ヘヨト ヘヨト



 $\sigma = t_1 t_3 t_4 t_2 t_3$ 

 $Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{0.4}$ 

÷.

ヘロト ヘヨト ヘヨト ヘヨト



 $\sigma = t_1 t_3 t_4 t_2 t_3$ 

 $Z_0 \xrightarrow{0.7} \stackrel{t_1}{\longrightarrow} \stackrel{0.0}{\longrightarrow} \stackrel{t_3}{\longrightarrow} \stackrel{0.4}{\longrightarrow} \stackrel{t_4}{\longrightarrow}$ 



æ

ヘロト ヘロト ヘビト ヘビト



 $\sigma = t_1 t_3 t_4 t_2 t_3$ 

 $Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{0.4} \xrightarrow{t_4} \xrightarrow{1.2}$ 

æ

ヘロト 人間 とくほ とくほ とう



 $\sigma = t_1 t_3 t_4 t_2 t_3$ 

 $Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{0.4} \xrightarrow{t_4} \xrightarrow{1.2}$ 

æ

ヘロト 人間 とくほ とくほ とう



 $\sigma = t_1 t_3 t_4 t_2 t_3$ 

 $Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{0.4} \xrightarrow{t_4} \xrightarrow{1.2} \xrightarrow{t_2} \rightarrow$ 



æ

ヘロト ヘロト ヘヨト ヘヨト



 $\sigma = t_1 t_3 t_4 t_2 t_3$ 

 $Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{0.4} \xrightarrow{t_4} \xrightarrow{1.2} \xrightarrow{t_2} \xrightarrow{0.5}$ 

2

ヘロト ヘロト ヘヨト ヘヨト



 $\sigma = t_1 t_3 t_4 t_2 t_3$ 

 $Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{0.4} \xrightarrow{t_4} \xrightarrow{1.2} \xrightarrow{t_2} \xrightarrow{0.5}$ 

2

ヘロト ヘロト ヘヨト ヘヨト



 $\sigma = t_1 t_3 t_4 t_2 t_3$ 

 $Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{0.4} \xrightarrow{t_4} \xrightarrow{1.2} \xrightarrow{t_2} \xrightarrow{0.5} \xrightarrow{t_3}$ 

2

ヘロト 人間 とくほ とくほ とう



 $\sigma = t_1 t_3 t_4 t_2 t_3$ 

 $Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{0.4} \xrightarrow{t_4} \xrightarrow{1.2} \xrightarrow{t_2} \xrightarrow{0.5} \xrightarrow{t_3} \xrightarrow{1.4} Z$ 

au = 0.7 0.0 0.4 1.2 0.5 1.4

Э

イロト イポト イヨト イヨト



 $\sigma = t_1 t_3 t_4 t_2 t_3$ 

 $\sigma(\tau) := \mathbf{Z}_0 \xrightarrow{\mathbf{0.7}} \xrightarrow{t_1} \xrightarrow{\mathbf{0.0}} \xrightarrow{t_3} \xrightarrow{\mathbf{0.4}} \xrightarrow{t_4} \xrightarrow{\mathbf{1.2}} \xrightarrow{t_2} \xrightarrow{\mathbf{0.5}} \xrightarrow{t_3} \xrightarrow{\mathbf{1.4}} \mathbf{Z}$ 

au = 0.7 0.0 0.4 1.2 0.5 1.4

イロト イポト イヨト イヨト 二日



#### $\sigma = t_1 t_3 t_4 t_2 t_3$

 $\sigma(\tau) := z_0 \xrightarrow{\mathbf{0.7}} \xrightarrow{t_1} \underbrace{\mathbf{0.0}}_{\tau} \xrightarrow{t_3} \underbrace{\mathbf{0.4}}_{\tau} \xrightarrow{t_4} \xrightarrow{\mathbf{1.2}} \xrightarrow{t_2} \underbrace{\mathbf{0.5}}_{\tau} \xrightarrow{t_3} \xrightarrow{\mathbf{1.4}} z$  $\tau = 0.7 \ 0.0 \ 0.4 \ 1.2 \ 0.5 \ 1.4$ 

э

イロト イポト イヨト イヨト



 $\sigma = t_1 t_3 t_4 t_2 t_3$ 

$$m_{\sigma} = (1, 2, 2, 1, 1)$$

æ

ヘロン 人間 とくほど 人間と



A B A B A
 A
 B
 A
 A
 B
 A
 A
 B
 A
 A
 B
 A
 A
 B
 A
 A
 B
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A

< ∃ >



$$B_{\sigma} = \left\{ \begin{array}{ll} 0 \leq x_{0}, & x_{0} \leq 2, \\ 0 \leq x_{1}, & x_{2} \leq 2, \\ 0 \leq x_{2}, & x_{3} \leq 2, \\ 1 \leq x_{3}, & x_{0} + x_{1} + x_{2} + x_{3} + x_{4} + x_{5} \leq 5, \\ 0 \leq x_{4}, & x_{4} + x_{5} \leq 2, \\ 0 \leq x_{5} \end{array} \right\}$$

19/44

E

<ロ> <同> <同> < 同> < 同>

The run  $\sigma(\tau)$  with



is feasible.

20 / 44

<ロ> < 回 > < 回 > < 回 > < 回 > .

$$(m_{\sigma},\begin{pmatrix}1.9\\1.4\\1.4\\1.4\\4.2\\\ddagger\\z_{0}\overset{\sigma(\tau)}{\not\equiv}z$$

21 / 44

æ

ヘロア 人間 アメヨア・



E

・ロト ・ 同 ト ・ ヨ ト ・ ヨ ト





E

・ロト ・ 同 ト ・ ヨ ト ・ ヨ ト
# **Example - Continuation**

#### The runs

$$\sigma(\tau_1^*) := \mathbf{Z}_0 \xrightarrow{\mathbf{1}} \xrightarrow{t_1} \underbrace{\mathbf{0}} \xrightarrow{\mathbf{t}_3} \xrightarrow{\mathbf{1}} \xrightarrow{\mathbf{t}_4} \xrightarrow{\mathbf{1}} \xrightarrow{\mathbf{t}_2} \underbrace{\mathbf{0}} \xrightarrow{\mathbf{t}_3} \xrightarrow{\mathbf{1}} \left\lfloor \mathbf{Z} \right\rfloor$$

#### and

$$\sigma(\tau_2^*) := z_0 \xrightarrow{\mathbf{1}} \xrightarrow{t_1} \underbrace{\mathbf{0}} \xrightarrow{t_3} \underbrace{\mathbf{0}} \xrightarrow{t_4} \underbrace{\mathbf{2}} \xrightarrow{t_2} \underbrace{\mathbf{0}} \xrightarrow{t_3} \underbrace{\mathbf{2}} \xrightarrow{\mathbf{2}} \left[ z \right]$$

#### are also feasible in $\mathcal{Z}$ .

Э

・ロト ・ 同ト ・ ヨト ・ ヨト

## **Example - Continuation**

#### The runs

$$\sigma(\tau_1^*) := Z_0 \xrightarrow{\mathbf{1}} \xrightarrow{t_1} \underbrace{\mathbf{0}} \xrightarrow{t_3} \xrightarrow{\mathbf{1}} \xrightarrow{t_4} \underbrace{\mathbf{1}} \xrightarrow{t_2} \underbrace{\mathbf{0}} \xrightarrow{t_3} \xrightarrow{\mathbf{1}} \left[ Z \right]$$
$$\sigma(\tau) = Z_0 \xrightarrow{\mathbf{0}} \xrightarrow{t_1} \underbrace{\mathbf{0}} \xrightarrow{\mathbf{0}} \xrightarrow{t_3} \underbrace{\mathbf{0}} \xrightarrow{\mathbf{0}} \xrightarrow{t_4} \underbrace{\mathbf{1}} \xrightarrow{\mathbf{2}} \xrightarrow{t_2} \underbrace{\mathbf{0}} \xrightarrow{\mathbf{0}} \xrightarrow{t_3} \underbrace{\mathbf{1}} \xrightarrow{\mathbf{4}} Z$$
$$\sigma(\tau_2^*) := Z_0 \xrightarrow{\mathbf{1}} \xrightarrow{t_1} \underbrace{\mathbf{0}} \xrightarrow{\mathbf{1}} \xrightarrow{\mathbf{1}} \xrightarrow{\mathbf{0}} \xrightarrow{\mathbf{1}} \xrightarrow{\mathbf{0}} \xrightarrow{\mathbf{1}} \xrightarrow{\mathbf{1}} \xrightarrow{\mathbf{2}} \xrightarrow{\mathbf{0}} \xrightarrow{\mathbf{1}} \xrightarrow{\mathbf{1}} \xrightarrow{\mathbf{2}} \left[ Z \right]$$

are also feasible in  $\mathcal{Z}$ .

E

#### Theorem 1:

Let  $\mathcal{Z}$  be a TPN and  $\sigma = t_1 \cdots t_n$  be a feasible transition sequence in  $\mathcal{Z}$  with a feasable run  $\sigma(\tau)$  of  $\sigma$  ( $\tau = \tau_0 \dots \tau_n$ ) i.e.

$$z_0 \xrightarrow{\tau_0} \xrightarrow{t_1} \cdots \xrightarrow{t_n} \xrightarrow{\tau_n} z_n = (m_n, h_n),$$

and all  $\tau_i \in \mathbb{R}_0^+$ . Then, there exists a further feasible run  $\sigma(\tau^*)$ ,  $\tau^* = \tau_0^* \dots \tau_n^*$  of  $\sigma$  with

$$z_0 \xrightarrow{\tau_0^*} \xrightarrow{t_1} \cdots \xrightarrow{t_n} \xrightarrow{\tau_n^*} z_n^* = (m_n^*, h_n^*).$$

such that

23 / 44

ヘロト ヘポト ヘヨト ヘヨト

#### Theorem 1 – Continuation:

$$egin{aligned} & z_0 & \stackrel{ au_0}{\longrightarrow} \stackrel{t_1}{\longrightarrow} \cdots \stackrel{t_n}{\longrightarrow} \stackrel{ au_n}{\longrightarrow} z_n = (m_n, h_n), \ au_i \in \mathbb{R}^+_0, \ & z_0 & \stackrel{ au_0^*}{\longrightarrow} \stackrel{ au_1^*}{\longrightarrow} \cdots \stackrel{ au_n}{\longrightarrow} \stackrel{ au_n^*}{\longrightarrow} z_n^* = (m_n^*, h_n^*) \end{aligned}$$

• For each  $i, 0 \le i \le n$  the time  $\tau_i^*$  is a natural number.

2 For each enabled transition *t* at marking  $m_n(=m_n^*)$  it holds:

$$\bullet h_n^*(t) = \lfloor h_n(t) \rfloor$$

$$\mathbf{2} \quad \sum_{i=1}^{n} \tau_i^* = \lfloor \sum_{i=1}^{n} \tau_i \rfloor$$

For each transition t ∈ T it holds:
 t is ready to fire in z<sub>n</sub> iff t is also ready to fire in [z<sub>n</sub>].



< ロ > < 同 > < 回 > < 回 > <

#### Theorem 1 – Continuation:

$$\begin{array}{cccc} z_0 & \xrightarrow{\tau_0} & \xrightarrow{t_1} & \cdots & \xrightarrow{t_n} & \xrightarrow{\tau_n} & z_n = (m_n, h_n), \ \tau_i \in \mathbb{R}_0^+. \\ z_0 & \xrightarrow{\tau_0^*} & \xrightarrow{t_1} & \cdots & \xrightarrow{t_n} & \xrightarrow{\tau_n^*} & z_n^* = (m_n^*, h_n^*), \ \tau_i^* \in \mathbb{N}. \end{array}$$

• For each  $i, 0 \le i \le n$  the time  $\tau_i^*$  is a natural number.

2 For each enabled transition *t* at marking  $m_n(=m_n^*)$  it holds:

$$\bullet h_n^*(t) = \lfloor h_n(t) \rfloor$$

$$\mathbf{2} \quad \sum_{i=1}^{n} \tau_i^* = \lfloor \sum_{i=1}^{n} \tau_i \rfloor$$

For each transition t ∈ T it holds:
 t is ready to fire in z<sub>n</sub> iff t is also ready to fire in [z<sub>n</sub>].



< ロ > < 同 > < 回 > < 回 > <

Let  $X_{\sigma} := \{x_0, x_1, \dots, x_n\}$  and  $\beta_0 : X_{\sigma} \longrightarrow \mathbb{R}_0^+$  with  $\beta_0(x_i) := \tau_i$  for each  $i \in \{0, \dots, n\}$ .

The successive construction of the assignment  $\beta^*$  from  $\beta_0$ .



< ロ > < 同 > < 回 > < 回 >

Let  $X_{\sigma} := \{x_0, x_1, \dots, x_n\}$  and  $\beta_0 : X_{\sigma} \longrightarrow \mathbb{R}_0^+$  with  $\beta_0(x_i) := \tau_i$  for each  $i \in \{0, \dots, n\}$ .

$\beta$	$\beta(x_0)$	$\beta(x_1)$	• • •	$\beta(\mathbf{x}_{n-i})$	$\beta(x_{n-(i-1)})$	• • •	$\beta(x_{n-1})$	$\beta(\mathbf{x}_n)$
$\beta_0$	r	r	• • •	r	r	• • •	r	r
$\beta_1$	r	r	• • •	r	r	•••	r	k

The successive construction of the assignment  $\beta^*$  from  $\beta_0$ .



Let  $X_{\sigma} := \{x_0, x_1, \dots, x_n\}$  and  $\beta_0 : X_{\sigma} \longrightarrow \mathbb{R}_0^+$  with  $\beta_0(x_i) := \tau_i$  for each  $i \in \{0, \dots, n\}$ .

$\beta$	$\beta(x_0)$	$\beta(x_1)$	•••	$\beta(x_{n-i})$	$\beta(x_{n-(i-1)})$	•••	$\beta(x_{n-1})$	$\beta(x_n)$
$\beta_0$	r	r	•••	r	r	• • •	r	r
$\beta_1$	r	r	•••	r	r	• • •	r	k
$\beta_2$	r	r		r	r		k	k

The successive construction of the assignment  $\beta^*$  from  $\beta_0$ .

Let  $X_{\sigma} := \{x_0, x_1, \dots, x_n\}$  and  $\beta_0 : X_{\sigma} \longrightarrow \mathbb{R}^+_0$  with  $\beta_0(x_i) := \tau_i$  for each  $i \in \{0, \dots, n\}$ .

β	$\beta(x_0)$	$\beta(x_1)$	•••	$\beta(x_{n-i})$	$\beta(x_{n-(i-1)})$	•••	$\beta(x_{n-1})$	$\beta(x_n)$
$\beta_0$	r	r	•••	r	r	•••	r	r
$\beta_1$	r	r	• • •	r	r	• • •	r	k
$\beta_2$	r	r		r	r	• • •	k	k
•			•			•		

The successive construction of the assignment  $\beta^*$  from  $\beta_0$ .



Let  $X_{\sigma} := \{x_0, x_1, \dots, x_n\}$  and  $\beta_0 : X_{\sigma} \longrightarrow \mathbb{R}^+_0$  with  $\beta_0(x_i) := \tau_i$  for each  $i \in \{0, \dots, n\}$ .

$\beta$	$\beta(x_0)$	$\beta(x_1)$	• • •	$\beta(x_{n-i})$	$\beta(x_{n-(i-1)})$	•••	$\beta(x_{n-1})$	$\beta(x_n)$
$\beta_0$	r	r	• • •	r	r	• • •	r	r
$\beta_1$	r	r	• • •	r	r	• • •	r	k
$\beta_2$	r	r		r	r		k	k
:			•			•		
•			-			-		
$\beta_i$	r	r	•••	r	k	•••	k	k

The successive construction of the assignment  $\beta^*$  from  $\beta_0$ .

Let  $X_{\sigma} := \{x_0, x_1, \dots, x_n\}$  and  $\beta_0 : X_{\sigma} \longrightarrow \mathbb{R}^+_0$  with  $\beta_0(x_i) := \tau_i$  for each  $i \in \{0, \dots, n\}$ .

β	$\beta(x_0)$	$\beta(x_1)$	• • •	$\beta(x_{n-i})$	$\beta(x_{n-(i-1)})$	• • •	$\beta(x_{n-1})$	$\beta(x_n)$
$\beta_0$	r	r	•••	r	r	•••	r	r
$\beta_1$	r	r	• • •	r	r	• • •	r	k
$\beta_2$	r	r	• • •	r	r	• • •	k	k
÷			:			-		
$\beta_i$	r	r		r	k		k	k
÷			÷			÷		

The successive construction of the assignment  $\beta^*$  from  $\beta_0$ .

Let  $X_{\sigma} := \{x_0, x_1, \dots, x_n\}$  and  $\beta_0 : X_{\sigma} \longrightarrow \mathbb{R}^+_0$  with  $\beta_0(x_i) := \tau_i$  for each  $i \in \{0, \dots, n\}$ .

β	$\beta(x_0)$	$\beta(x_1)$	• • •	$\beta(\mathbf{x}_{n-i})$	$\beta(x_{n-(i-1)})$	• • •	$\beta(x_{n-1})$	$\beta(x_n)$
$\beta_0$	r	r	• • •	r	r	• • •	r	r
$\beta_1$	r	r	• • •	r	r	•••	r	k
$\beta_2$	r	r	• • •	r	r	• • •	k	k
÷			÷			÷		
$\beta_i$	r	r		r	k		k	k
÷			÷			÷		
Bn	r	k		k	k		k	k

The successive construction of the assignment  $\beta^*$  from  $\beta_0$ .



Let  $X_{\sigma} := \{x_0, x_1, \dots, x_n\}$  and  $\beta_0 : X_{\sigma} \longrightarrow \mathbb{R}^+_0$  with  $\beta_0(x_i) := \tau_i$  for each  $i \in \{0, \dots, n\}$ .

β	$\beta(x_0)$	$\beta(x_1)$	• • •	$\beta(x_{n-i})$	$\beta(x_{n-(i-1)})$	• • •	$\beta(x_{n-1})$	$\beta(\mathbf{x}_n)$
$\beta_0$	r	r	• • •	r	r	• • •	r	r
$eta_1$	r	r	• • •	r	r	• • •	r	k
$\beta_2$	r	r	• • •	r	r	• • •	k	k
÷			:			÷		
•			•			•	1.	1.
$\beta_i$	r	r	•••	r	ĸ	•••	ĸ	K
:			:			:		
B	r	k	•	k	k		k	k
$\rho_n$	/	~	•••	n	n	•••	n	n
$\beta^* := \beta_{n+1}$	k	k		k	k		k	k

The successive construction of the assignment  $\beta^*$  from  $\beta_0$ .

25 / 44

For each inequality  $c \in B_{\sigma}$  let s(c) be the sum of variables.



Position of the real number  $[\![s(c)]\!]_{\beta_0}$  and the integers  $\lfloor [\![s(c)]\!]_{\beta_0} \rfloor - 1$ ,  $\lfloor [\![s(c)]\!]_{\beta_0} \rfloor$ ,  $\lceil [\![s(c)]\!]_{\beta_0} \rceil$  and  $\lceil [\![s(c)]\!]_{\beta_0} \rceil + 1$ .

	Time Petri Nets	Rounding of Runs
$Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{0.4}$	$\xrightarrow{t_4} \xrightarrow{1.2} \xrightarrow{t_2}$	$\xrightarrow{0.5} \xrightarrow{t_3} \xrightarrow{1.4} (m_{\sigma}, (1.9, 1.4, 1.4, 1.4, 4.2, \sharp))$

	Time Petri Nets	Rounding of Runs
$Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{0.4}$	$\xrightarrow{t_4} \xrightarrow{1.2} \xrightarrow{t_2}$	$\xrightarrow{0.5} \xrightarrow{t_3} \xrightarrow{1.4} (m_{\sigma}, (1.9, 1.4, 1.4, 1.4, 4.2, \sharp))$

<ロ> < 回> < 回> < 三> < 三> < 三</p>

Care Care

Time Petri Nets	Rounding of Runs
$Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{0.4} \xrightarrow{t_4} \xrightarrow{1.2} \xrightarrow{t_2}$	$\xrightarrow{0.5} \xrightarrow{t_3} \xrightarrow{1} (m_{\sigma}, (1.9, 1.4, 1.4, 1.4, 4.2, \sharp))$

<ロ> < 回> < 回> < 三> < 三> < 三</p>

Care Care

Time Petri Nets	Rounding of Runs
$Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{0.4} \xrightarrow{t_4} \xrightarrow{1.2} \xrightarrow{t_2} \xrightarrow{t_3} \xrightarrow{t_4} t$	$\xrightarrow{0.5} \xrightarrow{t_3} \xrightarrow{1} (m_{\sigma}, (1.5, 1.0, 1.0, 1.0, 3.8, \sharp))$

<ロ> < 回> < 回> < 三> < 三> < 三</p>

Care Care

Time Petri Nets	Rounding of Runs
$Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{0.4} \xrightarrow{t_4} \xrightarrow{1.2} \xrightarrow{t_2} \xrightarrow{t_2}$	$\xrightarrow{0.5} \xrightarrow{t_3} \xrightarrow{1} (m_{\sigma}, (1.5, 1.0, 1.0, 1.0, 3.8, \sharp))$

▲□▶▲圖▶▲≣▶▲≣▶ ■ のへで

	Time Petri Nets	Rounding of Runs
$Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{0.4}$	$\xrightarrow{t_4} \xrightarrow{1.2} \xrightarrow{t_2}$	$\xrightarrow{0} \xrightarrow{t_3} \xrightarrow{1} (m_{\sigma}, (1.5, 1.0, 1.0, 1.0, 3.8, \sharp))$

	Time Petri Nets	Rounding of Runs
$Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{0.4}$	$\xrightarrow{t_4} \xrightarrow{1.2} \xrightarrow{t_2}$	$\xrightarrow{0} \xrightarrow{t_3} \xrightarrow{1} (m_{\pi}, (1.0, 1.0, 1.0, 1.0, 3.3, \sharp))$

 $h_{\sigma}(t_1) = x_4 + x_5,$  $h_{\sigma}(t_5) = x_1 + x_2 + x_3 + x_4 + x_5$   $h_{\sigma}(t_2) = h_{\sigma}(t_3) = h_{\sigma}(t_4) = x_5$ 

◆□ > ◆□ > ◆臣 > ◆臣 > ─ 臣

Tim	e Petri Nets Ro	Rounding of Runs			
$Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{0.4} \xrightarrow{t_4}$	$\xrightarrow{1.2} \xrightarrow{t_2} \xrightarrow{0}$	$\xrightarrow{t_3} \xrightarrow{1} (m_{\sigma}, (1.0, 1.0, 1.0, 1.0, 1.0, 1.0, 1.0, 1.0, $	0,1.0,3.3,♯))		

$$h_{\sigma}(t_1) = x_4 + x_5, h_{\sigma}(t_5) = x_1 + x_2 + x_3 + x_4 + x_5$$

$$h_{\sigma}(t_2) = h_{\sigma}(t_3) = h_{\sigma}(t_4) = x_5$$

	Time Petri Nets	Rounding of Runs
$Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{0.4}$	$\xrightarrow{t_4} \xrightarrow{1} \xrightarrow{t_2}$	$\stackrel{0}{\longrightarrow} \stackrel{t_3}{\longrightarrow} \stackrel{1}{\longrightarrow} (m_{\sigma}, (1.0, 1.0, 1.0, 1.0, 3.3, \sharp))$

$$\begin{aligned} h_{\sigma}(t_1) &= x_4 + x_5, \\ h_{\sigma}(t_5) &= x_1 + x_2 + x_3 + x_4 + x_5 \end{aligned}$$

$$h_{\sigma}(t_2) = h_{\sigma}(t_3) = h_{\sigma}(t_4) = x_5$$

▲□▶▲□▶▲□▶▲□▶ ■ うへで

27 / 44

	Time Petri Nets	Rounding of Runs
$Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{0.4}$	$\xrightarrow{t_4} \xrightarrow{1} \xrightarrow{t_2}$	$\xrightarrow{0} \xrightarrow{t_3} \xrightarrow{1} (m_{\pi}, (1.0, 1.0, 1.0, 1.0, 3.1, \sharp))$

	I		<i>x</i> <sub>0</sub>	<i>x</i> <sub>1</sub>	<i>x</i> <sub>2</sub>	<i>x</i> 3	<i>x</i> <sub>4</sub>	<i>x</i> 5	$h_{\sigma}(t_1)$	$h_{\sigma}(t_2)$	$h_{\sigma}(t_5)$
$\hat{\beta}$	=	$\beta_0$	0.7	0.0	0.4	1.2	0.5	1.4	1.9	1.4	4.2
		$\beta_1$	0.7	0.0	0.4	1.2	0.5	1	1.5	1.0	3.8
		$\beta_2$	0.7	0.0	0.4	1.2	0	1	1.0		3.3
		$\beta_3$	0.7	0.0	0.4	1	0	1			3.1

$$\begin{aligned} h_{\sigma}(t_1) &= x_4 + x_5, \\ h_{\sigma}(t_5) &= x_1 + x_2 + x_3 + x_4 + x_5 \end{aligned}$$

 $h_{\sigma}(t_2) = h_{\sigma}(t_3) = h_{\sigma}(t_4) = x_5$ 

◆□→ ◆□→ ◆注→ ◆注→ □注

	Time Petri Nets	Rounding of Runs
$Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{0.4}$	$\xrightarrow{t_4} \xrightarrow{1} \xrightarrow{t_2}$	$\xrightarrow{0} \xrightarrow{t_3} \xrightarrow{1} (m_{\sigma}, (1.0, 1.0, 1.0, 1.0, 3.1, \sharp))$

$$h_{\sigma}(t_1) = x_4 + x_5, h_{\sigma}(t_5) = x_1 + x_2 + x_3 + x_4 + x_5$$

$$h_{\sigma}(t_2) = h_{\sigma}(t_3) = h_{\sigma}(t_4) = x_5$$

	Time Petri Nets	Rounding of Runs
$Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{1}$	$\xrightarrow{t_4} \xrightarrow{1} \xrightarrow{t_2}$	$\xrightarrow{0} \xrightarrow{t_3} \xrightarrow{1} (m_{\sigma}, (1.0, 1.0, 1.0, 1.0, 3.1, \sharp))$

$$h_{\sigma}(t_1) = x_4 + x_5, h_{\sigma}(t_5) = x_1 + x_2 + x_3 + x_4 + x_5$$

$$h_{\sigma}(t_2) = h_{\sigma}(t_3) = h_{\sigma}(t_4) = x_5$$

	Time Petri Nets	Rounding of Runs
$Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{1}$	$\xrightarrow{t_4} \xrightarrow{1} \xrightarrow{t_2}$	$\xrightarrow{0} \xrightarrow{t_3} \xrightarrow{1} (m_{\sigma}, (1.0, 1.0, 1.0, 1.0, 3.7, \sharp))$

$$h_{\sigma}(t_1) = x_4 + x_5,$$
  

$$h_{\sigma}(t_5) = x_1 + x_2 + x_3 + x_4 + x_5$$

$$h_{\sigma}(t_2) = h_{\sigma}(t_3) = h_{\sigma}(t_4) = x_5$$

▲□▶▲□▶▲□▶▲□▶ ■ うへで

27 / 44

	Time Petri Nets	Rounding of Runs
$Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{1}$	$\xrightarrow{t_4} \xrightarrow{1} \xrightarrow{t_2}$	$\xrightarrow{0} \xrightarrow{t_3} \xrightarrow{1} (m_{\sigma}, (1.0, 1.0, 1.0, 1.0, 3.7, \sharp))$

I		<i>x</i> <sub>0</sub>	<i>x</i> <sub>1</sub>	<i>x</i> <sub>2</sub>	<i>x</i> 3	<i>x</i> <sub>4</sub>	<i>x</i> 5	$  h_{\sigma}(t_1)$	$h_{\sigma}(t_2)$	$h_{\sigma}(t_5)$
$\hat{\beta} =$	$\beta_0$	0.7	0.0	0.4	1.2	0.5	1.4	1.9	1.4	4.2
	$\beta_1$	0.7	0.0	0.4	1.2	0.5	1	1.5	1.0	3.8
	$\beta_2$	0.7	0.0	0.4	1.2	0	1	1.0		3.3
	$\beta_3$	0.7	0.0	0.4	1	0	1			3.1
	$\beta_4$	0.7	0.0	1	1	0	1			3.7
	$\beta_5$	0.7		1	1	0	1			
	1							1		
L (L)						6	(1)	L (L)	I. (	

$$m_{\sigma}(t_1) = x_4 + x_5,$$
  
 $h_{\sigma}(t_5) = x_1 + x_2 + x_3 + x_4 + x_5$ 

$$h_{\sigma}(t_2) = h_{\sigma}(t_3) = h_{\sigma}(t_4) = x_5$$

▲□▶▲□▶▲□▶▲□▶ ■ うへで

27 / 44

	Time Petri Nets	Rounding of Runs
$Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0} \xrightarrow{t_3} \xrightarrow{1}$	$\xrightarrow{t_4} \xrightarrow{1} \xrightarrow{t_2}$	$\xrightarrow{0} \xrightarrow{t_3} \xrightarrow{1} (m_{\sigma}, (1.0, 1.0, 1.0, 1.0, 3.7, \sharp))$

I		<i>x</i> <sub>0</sub>	<i>x</i> <sub>1</sub>	<i>x</i> <sub>2</sub>	<i>x</i> 3	<i>x</i> <sub>4</sub>	<i>x</i> 5	$h_{\sigma}(t_1) h_{\sigma}(t_2) h_{\sigma}(t_5)$			
$\hat{\beta} =$	$\beta_0$	0.7	0.0	0.4	1.2	0.5	1.4	1.9	1.4	4.2	
	$\beta_1$	0.7	0.0	0.4	1.2	0.5	1	1.5	1.0	3.8	
	$\beta_2$	0.7	0.0	0.4	1.2	0	1	1.0		3.3	
	$\beta_3$	0.7	0.0	0.4	1	0	1			3.1	
	$\beta_4$	0.7	0.0	1	1	0	1			3.7	
	$\beta_5$	0.7	0	1	1	0	1				
$h_{\sigma}(t_1) = d$	<b>x</b> 4 +	- <b>X</b> 5,				$h_c$	$(t_2) =$	$= h_{\sigma}(t_3)$	$=h_{\sigma}($	$(t_4) = x_5$	

 $h_{\sigma}(t_5) = x_1 + x_2 + x_3 + x_4 + x_5$ 

<□> 
<□> 
<□> 
□>

	Time Petri Nets	Rounding of Runs
$Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0} \xrightarrow{t_3} \xrightarrow{1}$	$\xrightarrow{t_4} \xrightarrow{1} \xrightarrow{t_2}$	$\xrightarrow{0} \xrightarrow{t_3} \xrightarrow{1} (m_{\sigma}, (1.0, 1.0, 1.0, 1.0, 3.7, \sharp))$

I		<i>x</i> <sub>0</sub>	<i>x</i> <sub>1</sub>	<i>x</i> <sub>2</sub>	<i>x</i> <sub>3</sub>	<i>x</i> <sub>4</sub>	<i>x</i> 5	$h_{\sigma}(t_1) h_{\sigma}(t_2) h_{\sigma}(t_5)$		
$\hat{\beta} =$	$\beta_0$	0.7	0.0	0.4	1.2	0.5	1.4	1.9	1.4	4.2
	$\beta_1$	0.7	0.0	0.4	1.2	0.5	1	1.5	1.0	3.8
	$\beta_2$	0.7	0.0	0.4	1.2	0	1	1.0		3.3
	$\beta_3$	0.7	0.0	0.4	1	0	1			3.1
	$\beta_4$	0.7	0.0	1	1	0	1			3.7
	$\beta_5$	0.7	0	1	1	0	1			3.7
								•		

$$\begin{aligned} h_{\sigma}(t_1) &= x_4 + x_5, \\ h_{\sigma}(t_5) &= x_1 + x_2 + x_3 + x_4 + x_5 \end{aligned} \qquad h_{\sigma}(t_2) &= h_{\sigma}(t_3) = h_{\sigma}(t_4) = x_5 \end{aligned}$$

	Time Petri Nets	Rounding of Runs
$Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0} \xrightarrow{t_3} \xrightarrow{1} \xrightarrow{-}$	$\xrightarrow{t_4} \xrightarrow{1} \xrightarrow{t_2} \xrightarrow{-}$	$\xrightarrow{0} \xrightarrow{t_3} \xrightarrow{1} (m_{\sigma}, (1.0, 1.0, 1.0, 1.0, 3.7, \sharp))$

	I		<i>x</i> <sub>0</sub>	<i>x</i> <sub>1</sub>	<i>x</i> <sub>2</sub>	<i>x</i> 3	<i>x</i> <sub>4</sub>	<i>x</i> 5	$h_{\sigma}(t_1) h_{\sigma}(t_2) h_{\sigma}(t_5)$		
$\hat{\beta}$	=	$\beta_0$	0.7	0.0	0.4	1.2	0.5	1.4	1.9	1.4	4.2
		$\beta_1$	0.7	0.0	0.4	1.2	0.5	1	1.5	1.0	3.8
		$\beta_2$	0.7	0.0	0.4	1.2	0	1	1.0		3.3
		$\beta_3$	0.7	0.0	0.4	1	0	1			3.1
		$\beta_4$	0.7	0.0	1	1	0	1			3.7
		$\beta_5$	0.7	0	1	1	0	1			3.7
$\beta^*$	=	$\beta_{6}$		0	1	1	0	1			

 $\begin{aligned} h_{\sigma}(t_1) &= x_4 + x_5, \\ h_{\sigma}(t_5) &= x_1 + x_2 + x_3 + x_4 + x_5 \end{aligned}$ 

$$h_{\sigma}(t_2) = h_{\sigma}(t_3) = h_{\sigma}(t_4) = x_5$$

< □ > < □ > < □ > < □ > < □ > = □

	Time Petri Nets	Rounding of Runs
$Z_0 \xrightarrow{1} \xrightarrow{t_1} \xrightarrow{0} \xrightarrow{t_3} \xrightarrow{1}$	$\xrightarrow{t_4} \xrightarrow{1} \xrightarrow{t_2}$	$\xrightarrow{0} \xrightarrow{t_3} \xrightarrow{1} (m_{\sigma}, (1.0, 1.0, 1.0, 1.0, 3.7, \sharp))$

	Т		<i>x</i> <sub>0</sub>	<i>x</i> <sub>1</sub>	<i>x</i> <sub>2</sub>	<i>x</i> 3	<i>x</i> <sub>4</sub>	<i>x</i> 5	$h_{\sigma}(t_1)$	$h_{\sigma}(t_2)$	$h_{\sigma}(t_5)$
$\hat{\beta}$	=	$\beta_0$	0.7	0.0	0.4	1.2	0.5	1.4	1.9	1.4	4.2
		$\beta_1$	0.7	0.0	0.4	1.2	0.5	1	1.5	1.0	3.8
		$\beta_2$	0.7	0.0	0.4	1.2	0	1	1.0		3.3
		$\beta_3$	0.7	0.0	0.4	1	0	1			3.1
		$\beta_4$	0.7	0.0	1	1	0	1			3.7
		$\beta_5$	0.7	0	1	1	0	1			3.7
$\beta^*$	=	$\beta_{6}$	1	0	1	1	0	1			

 $\begin{aligned} h_{\sigma}(t_1) &= x_4 + x_5, \\ h_{\sigma}(t_5) &= x_1 + x_2 + x_3 + x_4 + x_5 \end{aligned}$ 

$$h_{\sigma}(t_2) = h_{\sigma}(t_3) = h_{\sigma}(t_4) = x_5$$

< □ > < □ > < □ > < □ > < □ > = □

	Time Petri Nets	Rounding of Runs
$Z_0 \xrightarrow{1} \xrightarrow{t_1} \xrightarrow{0} \xrightarrow{t_3} \xrightarrow{1} \xrightarrow{t_4}$	$t_4 \rightarrow \frac{1}{\longrightarrow} \xrightarrow{t_2} -$	$\xrightarrow{0} \xrightarrow{t_3} \xrightarrow{1} (m_{\sigma}, (1.0, 1.0, 1.0, 1.0, 4.0, \sharp))$

	I		<i>x</i> <sub>0</sub>	<i>x</i> <sub>1</sub>	<i>x</i> <sub>2</sub>	<i>x</i> 3	<i>x</i> <sub>4</sub>	<i>x</i> 5	$h_{\sigma}(t_1)$	$h_{\sigma}(t_2)$	$h_{\sigma}(t_5)$
$\hat{\beta}$	=	$\beta_0$	0.7	0.0	0.4	1.2	0.5	1.4	1.9	1.4	4.2
		$\beta_1$	0.7	0.0	0.4	1.2	0.5	1	1.5	1.0	3.8
		$\beta_2$	0.7	0.0	0.4	1.2	0	1	1.0		3.3
		$\beta_3$	0.7	0.0	0.4	1	0	1			3.1
		$\beta_4$	0.7	0.0	1	1	0	1			3.7
		$\beta_5$	0.7	0	1	1	0	1			3.7
$\beta^*$	=	$\beta_{6}$	1	0	1	1	0	1			4.0

 $h_{\sigma}(t_1) = x_4 + x_5,$  $h_{\sigma}(t_5) = x_1 + x_2 + x_3 + x_4 + x_5$   $h_{\sigma}(t_2) = h_{\sigma}(t_3) = h_{\sigma}(t_4) = x_5$ 

(ロ) (部) (E) (E) (E)

#### Theorem 2:

Let  $\mathcal{Z}$  be a TPN and  $\sigma = t_1 \cdots t_n$ ) be a feasible transition sequence in  $\mathcal{Z}$ , with feasable run  $\sigma(\tau)$  of  $\sigma$  ( $\tau = \tau_0 \dots \tau_n$ ) i.e.

$$z_0 \xrightarrow{\tau_0} \xrightarrow{t_1} \cdots \xrightarrow{t_n} \xrightarrow{\tau_n} z_n = (m_n, h_n),$$

and all  $\tau_i \in \mathbb{R}^+_0$ . Then, there exists a further feasible run  $\sigma(\tau^*)$  of  $\sigma$  with

$$z_0 \xrightarrow{\tau_0^*} \xrightarrow{t_1} \cdots \xrightarrow{t_n} \xrightarrow{\tau_n^*} z_n^* = (m_n^*, h_n^*).$$

such that

#### Theorem 2 – Continuation:

- For each  $i, 0 \le i \le n$  the time  $\tau_i^*$  is a natural number.
- **2** For each enabled transition *t* at marking  $m_n(=m_n^*)$  it holds:

$$h_n(t)^* = \lceil h_n(t) \rceil.$$

$$2 \quad \sum_{i=1}^{n} \tau_i^* = \left\lceil \sum_{i=1}^{n} \tau_i \right\rceil$$

For each transition t ∈ T holds:
 t is ready to fire in z<sub>n</sub> iff t is also ready to fire in [z<sub>n</sub>].

< ロ > < 同 > < 回 > < 回 >

	Time Petri Nets	Rounding of Runs
$Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{0.4}$	$\xrightarrow{t_4} \xrightarrow{1.2} \xrightarrow{t_2}$	$\xrightarrow{0.5} \xrightarrow{t_3} \xrightarrow{1.4} (m_{\sigma}, (1.9, 1.4, 1.4, 1.4, 4.2, \sharp))$

	Ш		x <sub>0</sub>	<i>x</i> <sub>1</sub>	<i>x</i> <sub>2</sub>	<i>x</i> 3	<i>x</i> <sub>4</sub>	<i>x</i> 5	$h_{\sigma}(t_1)$	$h_{\sigma}(t_2)$	$h_{\sigma}(t_5)$
$\hat{\beta}$	=	$\beta_0$	0.7	0.0	0.4	1.2	0.5	1.4	1.9	1.4	4.2
		$\beta_1$	0.7	0.0	0.4	1.2	0.5	2	2.5	2.0	4.8
		$\beta_2$	0.7	0.0	0.4	1.2	0	2	2.0		4.3
		$\beta_3$	0.7	0.0	0.4	2	0	2			5.1
		$\beta_4$	0.7	0.0	0	2	0	2			4.7
		$\beta_5$	0.7	0	0	2	0	2			4.7
		$\beta_{6}$	1	0	0	2	0	2			5.0

 $z_{0} \xrightarrow{1} \xrightarrow{t_{1}} \xrightarrow{0} \xrightarrow{t_{3}} \xrightarrow{0} \xrightarrow{t_{4}} \xrightarrow{2} \xrightarrow{t_{2}} \xrightarrow{0} \xrightarrow{t_{3}} \xrightarrow{2} (m_{\sigma}, (2, 2, 2, 2, 5, \sharp))$ 

< □ > < □ > < □ > < □ > < □ > = □
	Time Petri Nets	Rounding of Runs
$Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{0.4}$	$\xrightarrow{t_4} \xrightarrow{1.2} \xrightarrow{t_2}$	$\xrightarrow{0.5} \xrightarrow{t_3} \xrightarrow{1.4} (m_{\pi}, (1.9, 1.4, 1.4, 1.4, 4.2, \sharp))$

### The time length of the run $\sigma(\tau)$ is $\hat{\beta}(x_0) + \hat{\beta}(x_1) + \hat{\beta}(x_2) + \hat{\beta}(x_3) + \hat{\beta}(x_4) + \hat{\beta}(x_5) = 4.2$



Time Petri No	Nets Rounding of Runs	
$Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{0.4} \xrightarrow{t_4} \xrightarrow{1.2} \xrightarrow{-1}$	$\xrightarrow{t_2} \xrightarrow{0.5} \xrightarrow{t_3} \xrightarrow{1.4} (m_{\sigma}, (1.9, 1.4, 1.4, 1.4, 4.1)$	2, ♯))

The time length of the run  $\sigma(\tau)$  is  $\hat{\beta}(x_0) + \hat{\beta}(x_1) + \hat{\beta}(x_2) + \hat{\beta}(x_3) + \hat{\beta}(x_4) + \hat{\beta}(x_5) = 4.2$ In tableau I: The time length of the run  $\sigma(\tau_1^*)$  is 4



	Time Petri Nets	Rounding of Runs
$Z_0 \xrightarrow{0.7} \xrightarrow{t_1} \xrightarrow{0.0} \xrightarrow{t_3} \xrightarrow{0.4}$	$\xrightarrow{t_4} \xrightarrow{1.2} \xrightarrow{t_2}$	$\xrightarrow{0.5} \xrightarrow{t_3} \xrightarrow{1.4} (m_{\sigma}, (1.9, 1.4, 1.4, 1.4, 4.2, \sharp))$

The time length of the run  $\sigma(\tau)$  is  $\hat{\beta}(x_0) + \hat{\beta}(x_1) + \hat{\beta}(x_2) + \hat{\beta}(x_3) + \hat{\beta}(x_4) + \hat{\beta}(x_5) = 4.2$ In tableau I: The time length of the run  $\sigma(\tau_1^*)$  is 4 In tableau II: The time length of the run  $\sigma(\tau_2^*)$  is 5



イロト イポト イヨト イヨト

#### Some Conclusions

- Each feasible transitions sequence σ in Z can be realized with an integer run.
- Each reachable *p*-marking in  $\mathcal{Z}$  can be reached using **integer** runs only.
- If z is reachable in  $\mathcal{Z}$ , then  $\lfloor z \rfloor$  and  $\lceil z \rceil$  are reachable in  $\mathcal{Z}$  as well.
- The length of the shortest and longest time path (if this is finite) between two arbitrary *p*-markings are natural numbers.

A run  $\sigma(\tau) = \tau_0 \ t_1 \ \tau_1 \dots t_n \ \tau_n$  is an **integer** one, if  $\tau_i \in \mathbb{N}$  for each  $i = 0 \dots n$ .



#### Integer States

A state z = (m, h) is an **integer** one, if  $h(t) \in \mathbb{N}$  for each in *m* enabled transition *t*.

#### Theorem 3:

Let  $\mathcal{Z}$  be a finite TPN, i.e.  $lft(t) \neq \infty$  for all  $t \in T$ . The set of all reachable integer states in  $\mathcal{Z}$  is finite if and only if the set of all reachable *p*-markings in  $\mathcal{Z}$  is finite.

#### Integer States

A state z = (m, h) is an **integer** one, if  $h(t) \in \mathbb{N}$  for each in *m* enabled transition *t*.

#### Theorem 3:

Let  $\mathcal{Z}$  be a finite TPN, i.e.  $lft(t) \neq \infty$  for all  $t \in T$ . The set of all reachable integer states in  $\mathcal{Z}$  is finite if and only if the set of all reachable *p*-markings in  $\mathcal{Z}$  is finite.

#### Remark:

Theorem 3 can be generalized for all TPNs (applying a further reduction of the state space).

イロト イポト イヨト イヨト

#### Modified Rule

Let  $\mathcal{Z}$  be an arbitrary TPN. The state change **by time elapsing** can be slightly **modified** for each transition *t* with  $lft(t) = \infty$ , because in order to fire such a transition *t* 

- it is important to know whether t is old enough or not, i.e. whether t has been enabled last for eft(t) (or more) time units or t is younger.
- Thus, the time h(t) increases until eft(t). After that,
  the clock of t remains in this position (although the time is elapsing), unless t becomes disabled.



#### Theorem 4:

In an arbitrary TPN a *p*-marking is reachable using the nonmodified definition iff it is reachable using the modified one.



#### Theorem 4:

In an arbitrary TPN a *p*-marking is reachable using the nonmodified definition iff it is reachable using the modified one.

All **integer states** in an arbitrary TPN, obtained by using the **modified definition**, are called the **essential states** of this net.



くロト く伺 とくき とくきつ

#### Theorem 4:

In an arbitrary TPN a *p*-marking is reachable using the nonmodified definition iff it is reachable using the modified one.

All **integer states** in an arbitrary TPN, obtained by using the **modified definition**, are called the **essential states** of this net.

#### **Theorem 5:**

An arbitrary TPN is bounded iff the set of its reachable essential states is finite.

< ロ > < 同 > < 回 > < 国

#### Standard Rule vs. Modified Rule



It holds that:

$$RIS_{\mathcal{Z}} := \{ (1, \begin{pmatrix} 0\\0 \end{pmatrix}), (1, \begin{pmatrix} 1\\1 \end{pmatrix}), (1, \begin{pmatrix} 2\\2 \end{pmatrix}), (1, \begin{pmatrix} 3\\3 \end{pmatrix}), (1, \begin{pmatrix} 4\\4 \end{pmatrix}), (1, \begin{pmatrix} 5\\5 \end{pmatrix}), (0, \begin{pmatrix} \sharp\\ \sharp \end{pmatrix}) \}$$
  
and  
$$REIS_{\mathcal{Z}} := \{ (1, \begin{pmatrix} 0\\0 \end{pmatrix}), (1, \begin{pmatrix} 1\\1 \end{pmatrix}), (1, \begin{pmatrix} 2\\2 \end{pmatrix}), (1, \begin{pmatrix} 3\\2 \end{pmatrix}), (1, \begin{pmatrix} 4\\2 \end{pmatrix}), (1, \begin{pmatrix} 5\\2 \end{pmatrix}), (0, \begin{pmatrix} \sharp\\ \sharp \end{pmatrix}) \}.$$

Clearly, neither  $REIS_{\mathcal{Z}} \subseteq RIS_{\mathcal{Z}}$  nor  $REIS_{\mathcal{Z}} \supseteq RIS_{\mathcal{Z}}$ .



・ コ ト ・ 雪 ト ・ 日 ト ・

#### Discrete Reduction of the State Space



The set of all reachable states



#### Discrete Reduction of the State Space



The set of all reachable states

The set of all essential states



 $\bullet \equiv \bullet$ 





E

・ロン ・回 と ・ 回 と ・ 回 と



The reachability graph is a weighted directed graph, including the time explicit.

<ロト < 同ト < 回ト <



The reachability graph is a weighted directed graph, including the time explicit.

<ロト < 同ト < 回ト <



The reachability graph is a weighted directed graph, including the time explicit.

< ロ > < 同 > < 三 >



The reachability graph is a weighted directed graph, including the time explicit.

< ロ > < 同 > < 回 > < 国



The reachability graph is a weighted directed graph, including the time explicit.

<ロト < 同ト < 回ト <



The reachability graph is a weighted directed graph, including the time explicit.



The reachability graph is a weighted directed graph, including the time explicit.

**Reachable Graph** 

## Example: A finite TPN and its reachability graph





## Example: A finite TPN and its reachability graph





Image: A math a math

## Example: A finite TPN and its reachability graph





э

イロト イポト イヨト イヨト

**Reachable Graph** 

## Example: A finite TPN and its reachability graph







э

イロト イポト イヨト イヨト

**Time Petri Nets** 

**Reachable Graph** 

# Example: A non-finite TPN and its reachability graph





э

< D > < P > < P > < P >

-∢ ≣ →

## State Classes Reachability Graphs vs. Essential States Reachability Graphs



ъ

< A >

Time Petri Nets

**Reachable Graph** 

number of	essenti	essential-states algorithm			state class algorithm		
tokens	number of	number of	total	number of	number of	total	
in <i>p</i> 0	vertices	edges	number	vertices	edges	number	
0	1	0	1	1	0	1	
1	4	21	25	4	5	9	
2	63	310	373	81	157	238	
3	250	1252	1502	258	574	832	
4	692	3920	4612	1053	2979	4032	
5	1367	8115	9482	2653	8119	10772	
6	2265	13769	16034	5000	15884	20884	
7	3386	20882	24268	8089	26315	34404	
8	4730	29454	34184	11909	39371	51280	
9	6297	39485	45782	16454	55023	71477	
10	8087	50975	59062	21708	73210	94918	

The firing rule is defined based on static conflict.



Э

・ロト ・聞 ト ・ ヨ ト ・ ヨ ト

Time Petri Nets

**Reachable Graph** 

number of	essential-states algorithm			state class algorithm		
tokens	number of	number of	total	number of	number of	total
in <i>p</i> 0	vertices	edges	number	vertices	edges	number
0	1	0	1	1	0	1
1	4	21	25	4	5	9
2	86	441	527	94	186	280
3	550	2740	3290	570	1354	1924
4	1916	9975	11891	2181	5907	8088
5	9167	50618	59785	16588	53781	70369
7	15152	84449	99601	34118	114249	148367
8	22862	127989	150851	61123	208195	269318
9	32165	180510	212675	97479	335218	432697
10	42989	241713	284702	142712	493602	636314

The firing rule is defined based on dynamic conflict.



43/44

Э

ヘロア 人間 アメヨア 人口 ア

## **Questions, Discussions ?**

