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CS&P Workshop 2021 Attack trees with Time Constraints

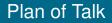
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A real-time attack on systems require a realtime measure to stop (or reduce the damage) the attack.

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Attack Trees

Challenges

Time constraints

Attack trees to timed automata

Conclusion





Graphical representation of ontology that :





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- classify i.e., shows the breakdown of threats and how they can be achieve





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- capture i.e., identifying root cause of threats (attacks)
- classify i.e., shows the breakdown of threats and how they can be achieve
- share i.e., study how to repel the attacks.



An attack tree is constructed in a bottom-up manner where values at the leaves of the attack will be propagated towards the root node of the tree.

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- Root node represent the attack target
- internal nodes (child nodes) are conditions which must be satisfied to reach the root, and
- leaf nodes are atomic actions that need to be carried out to initiate the attack
- nodes (except for leaves) are associated with gates refinement e.g AND and OR gates (can be easily extended to show sequential actions SAND,SOR)

Attack tree example 1



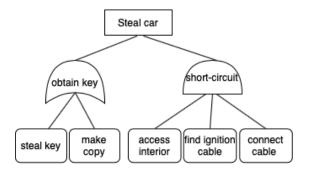


Figure: A simple example of attack tree



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- study and identify the possible attack paths that can be easily exploited [S. Haque et. al.]
- attack-defence and countermeasures to prevent attack from happening [Kordy, B et. al.]

Dynamic threat environment





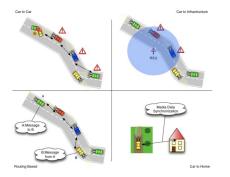
Dynamic threat environment



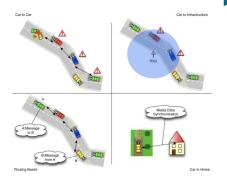


Most assets are no longer stand-along i.e., interacts with other set of systems and objects from the physical environment, e.g CPS.



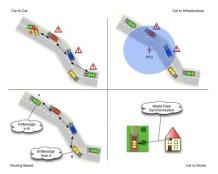






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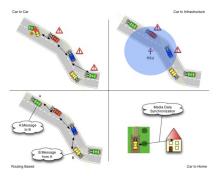




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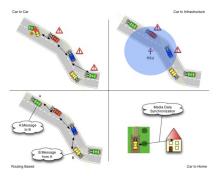
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Proposed approach



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Let T_r be an attack tree, we defined an attack as as a mapping $\mathcal{A} : \mathcal{Q}_L \rightarrow (Act \times T) \cup \{Nil\}$ where

- Q_L is a set of leaf nodes
- Act is a set of attack actions
- ► *T* is a set of attack time

We say that A is a simple attack if $A(I, t) \neq NiI$ only for one leaf node.



Let τ be a set of time constraints that is represented by a constant pair $\langle b, f \rangle$, with *b* marking the start of an attack and *f* marking the (expected) end of attack on the parent node such that $b, f \in \mathbb{Q}$ and $b \leq f$.

A set of gates $\mathcal{G} : \{q_0\} \cup \mathcal{Q}_S \rightarrow \{AND, OR\} \times \tau$.



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Given an attack tree T_r , a set of attack time T, and a set of attack actions *Act*. An attack over T_r can only succeed;



- ▶ for AND node if for all simple attacks A₁...A_n such that t₁...t_n are less than or equal to the constant time interval on the node,
- for OR node if there exists a simple attack such that its attack time is less than or equal to the constant time interval on the node.



► Let *T_r* be an attack tree and *τ* be a set of time constraints, we proposed how to formally define some properties of the tree via weighted timed automata.

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- ▶ this properties can be checked and verified using UPPAAL.
- We translated the tree into a parallel composition of weighted timed automata (WTA).

UPPAAL translation



- The sets of nodes of the attack trees are translated to a set of locations in the weighted timed automata.
- For each leaf node in the attack tree, we have a WTA that represents a linear path from the leaf to the root node.
- Altogether there are as many WTAs as the leaf nodes in the attack tree. Each location that represents a leaf which has a clock that is activated when there is an attack in the corresponding leaf in the tree.
- An attack on a node in the tree represents an enabled transition in the WTA. Initially, clocks become active when events synchronized, and end with either a success or fail synchronization action.



- If a location that represent a the sub_goal of the attack tree has an OR gate, a single event (i.e. action) together with a clock is activated and a transition to the target location is enabled
- otherwise all the events (i.e., AND gate) needs to be activated, and the target location can only be reached if all elapsing time in the clocks is less than (predefined) max (that represent the constant time intervals).



- We discussed how attack trees can be extended with a set of attack time, and a set of time constraints at the gates refinement
- we show how this extensions can be modelled and analyse using WTA
- We plan to study how actions observations, which cannot be identified as vulnerabilities with attack trees can be prevented from aiding an attack.

Thank you!